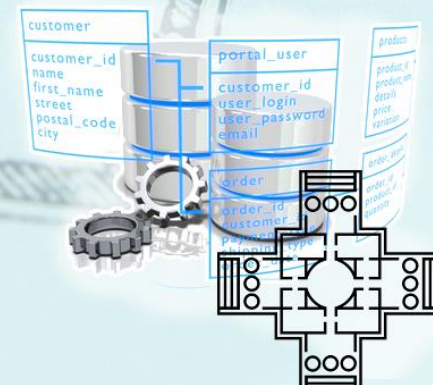


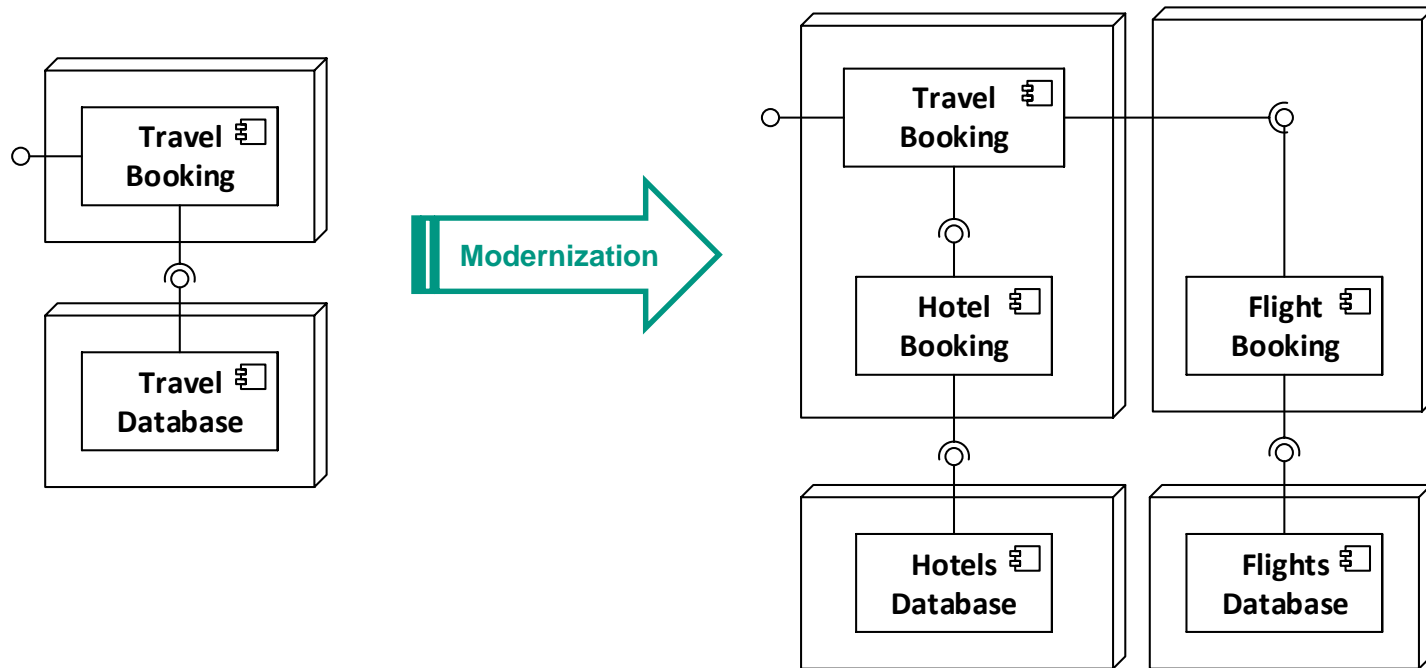
# Extending the Palladio Component Model to Analyze Data Contention for Modernizing Transactional Software Towards Service-Orientation

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# Modernizing towards Service-Orientation



Monolithic application

Component-based, service-oriented application

# Different Ways to Isolate Concurrent Transactions

- Isolation prevents concurrent transactions from mutually overriding their data

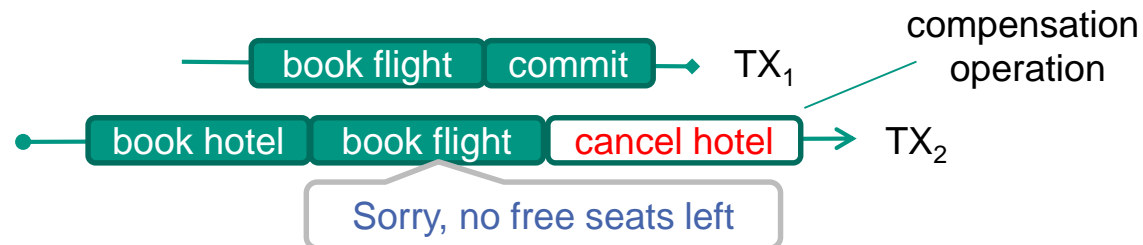
## Classical „pessimistic“ way

- Lock data before reading/writing
- Waiting time impacts performance

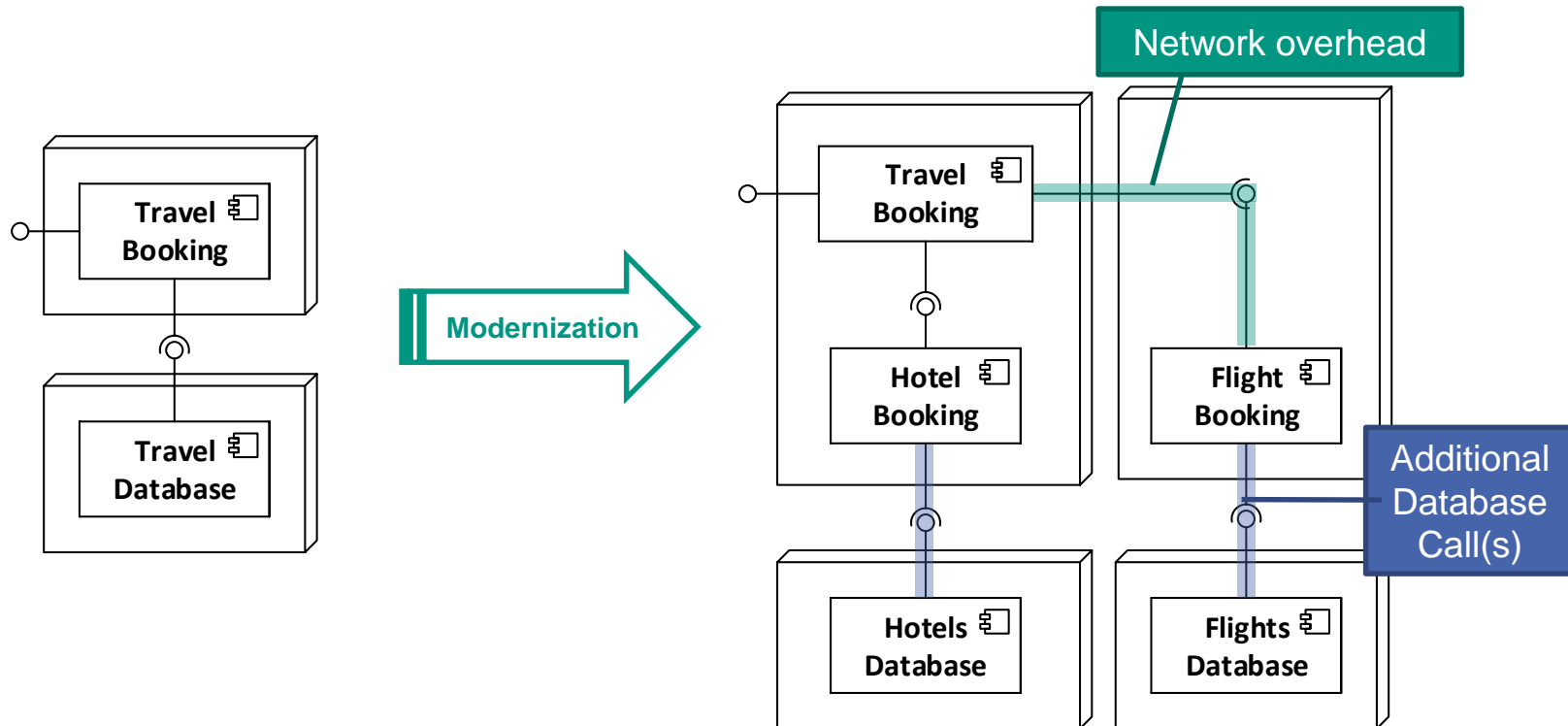


## Service-oriented „optimistic“ way

- Pretend there is no concurrent data access...
- ...until access conflict actually happens → abort + retry
- Abort often uses **compensation**
- Try-Cancel-Confirm

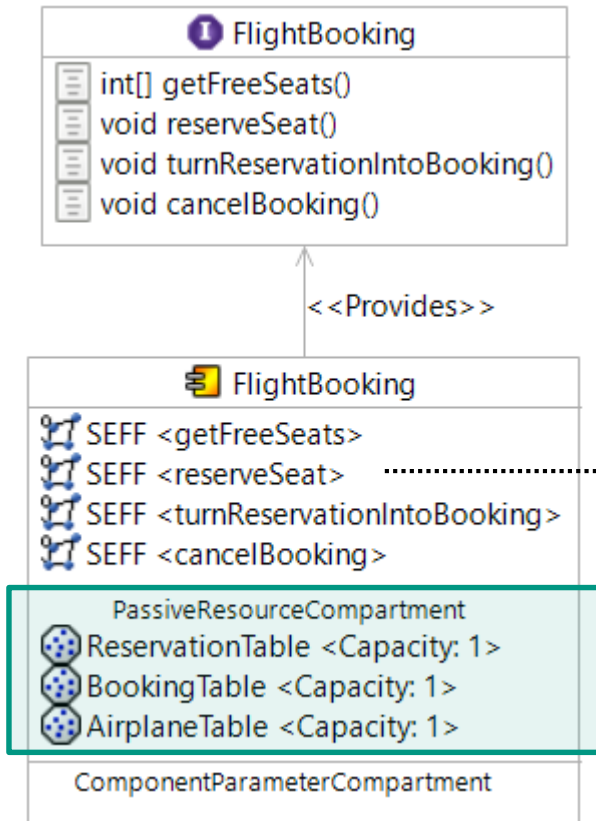


# Modernization Revisited

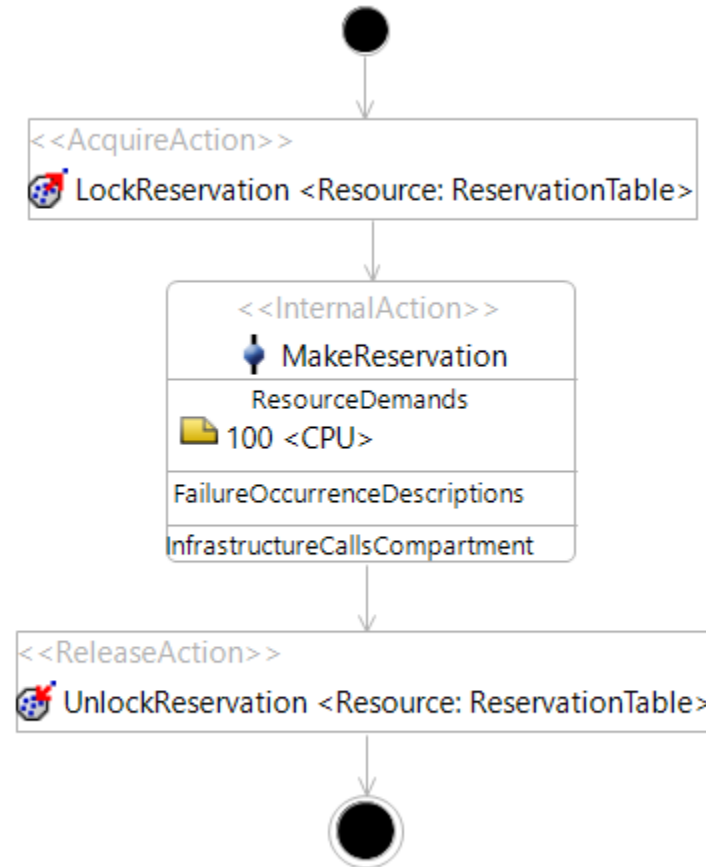


- Transaction duration increased by **network** + **database** overheads
    - Locks held longer (Pessimistic)
    - Higher abort risk (Optimistic)
- } Increased Data contention  
**Impact on system performance? → PCM**

# Transactions with the PCM – Status Quo



## RDSEFF *reserveSeat*:



# Transactions with the PCM – Status Quo

RDSEFF *reserveSeat*.

1 FlightBooking

```
int[] getFreeSeats()
void reserveSeat()
void turnReservationIntoBooking()
void turnBookingIntoReservation()
```

## Why **not** to use Passive Resources for Transaction Modeling:

- Modeler needs to be expert with locking schemes
- Chosen locking scheme scattered over multiple RDSEFFs
- Locking granularity: row or page level impractical
- Locking mode: shared mode impractical
- ...

```
ReservationTable <Capacity: 1>
Booking <Capacity: 1>
AirplaneTable <Capacity: 1>
```

ComponentParameterCompartment

UnlockReservation <Resource: ReservationTable>

# PCM.TX Overview



- **Transaction** Boundaries
- Compensation Behaviour
- „Try-Cancel-Confirm“ Behaviour
  
- **Queries**
- (Transaction Boundaries)
- (Compensation Behaviour)
  
- **Schema** Definition
- Table Deployment

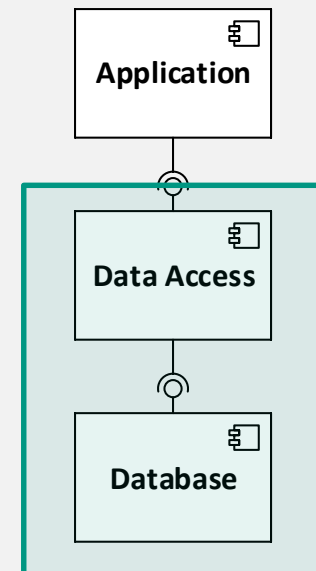
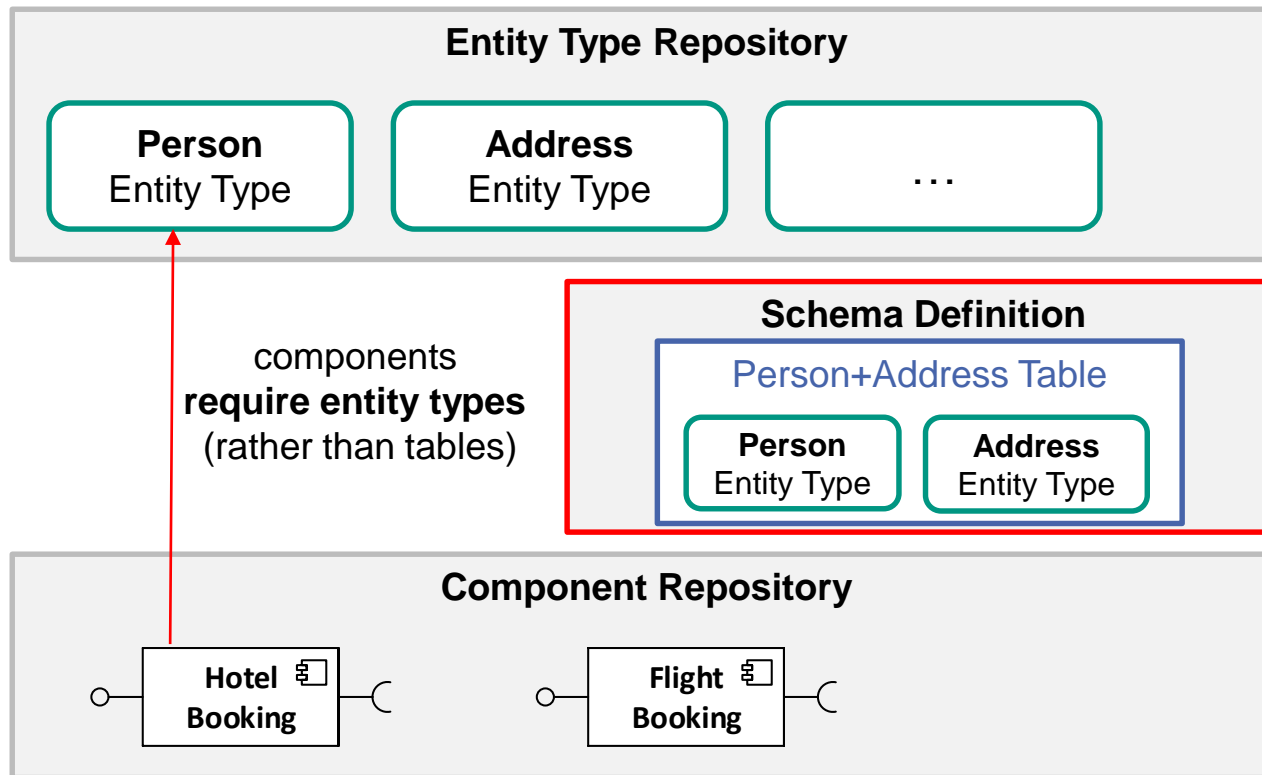
## Cross-Cutting:

Explicit Context (Database) Dependencies

Separate Schema Specification from  
Component Specification

# Separate Schema Specification from Component Specification

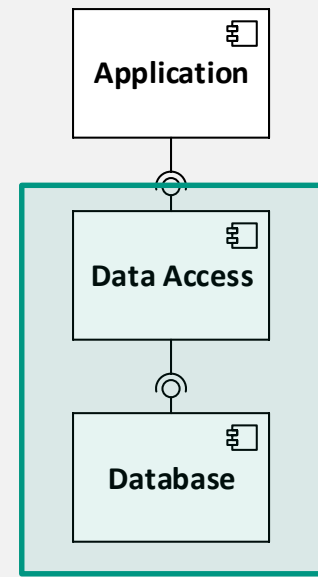
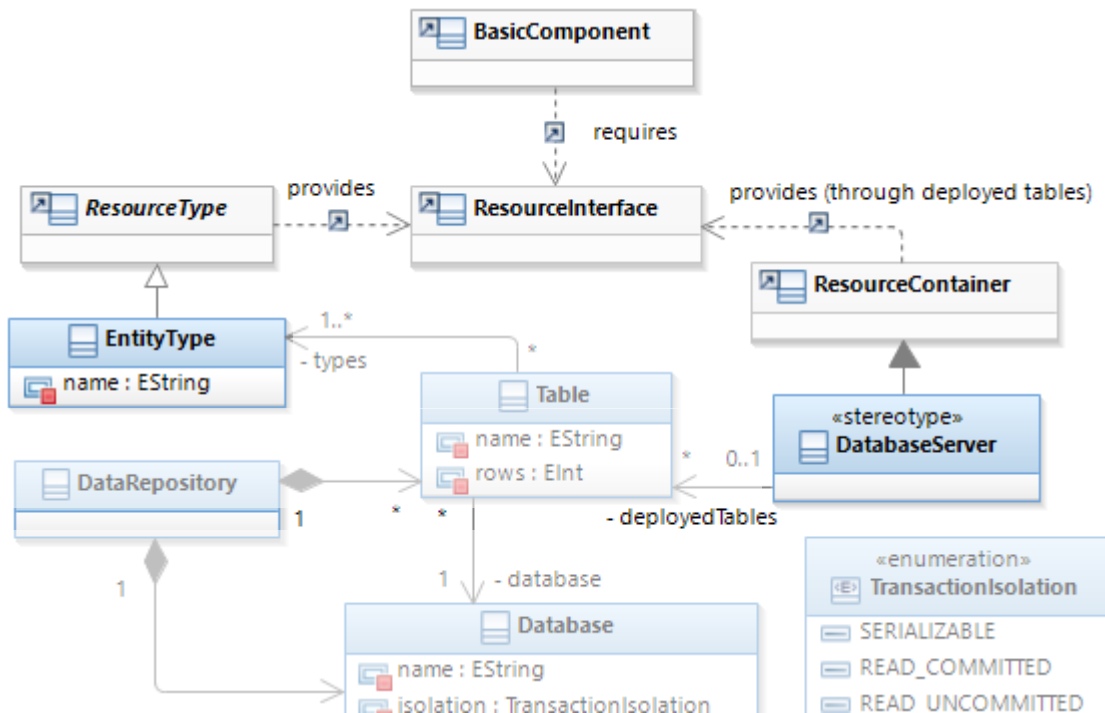
- Problem: If components refer to tables directly, they need modification whenever the schema changes!
- Solution: **Decouple** schema and component specification



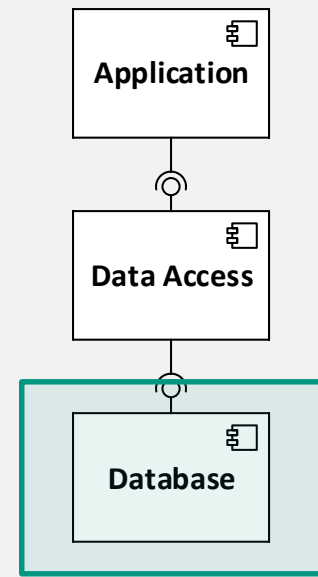
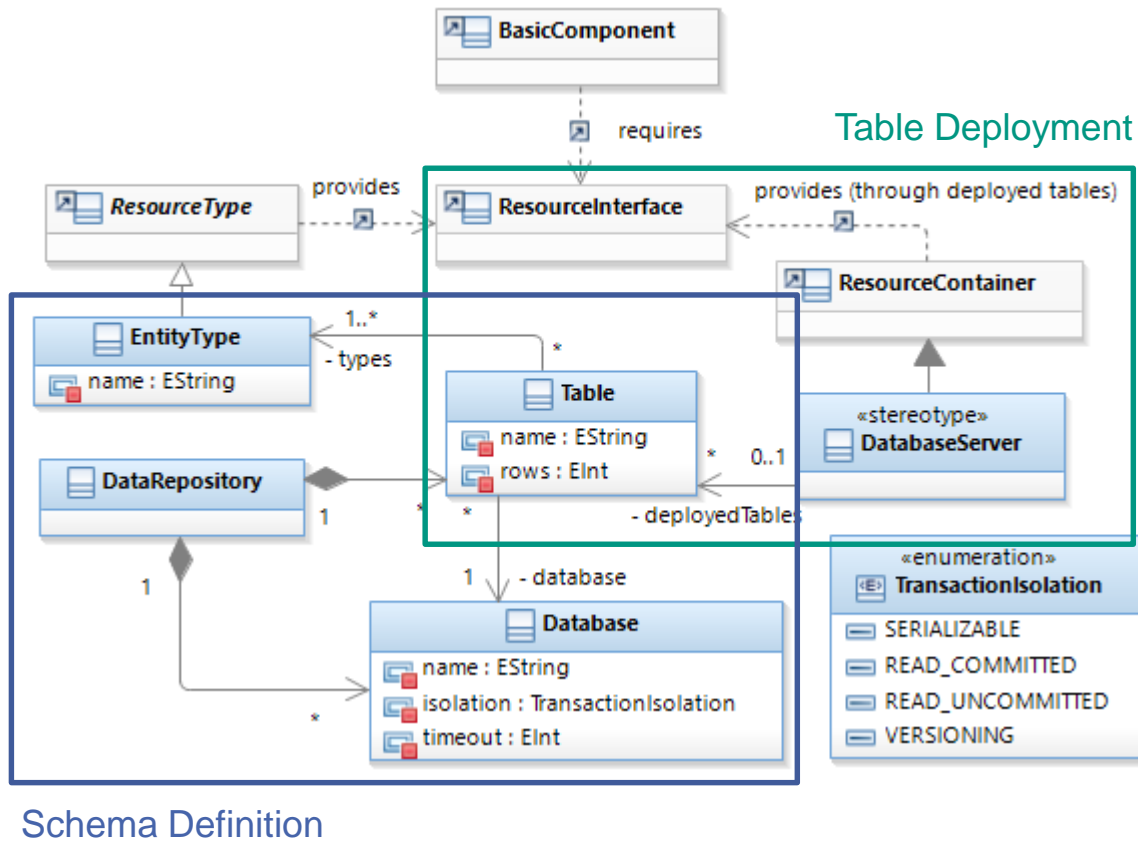


# Explicit Context (Database) Dependencies

- “A software component is a unit of composition with contractually specified interfaces and explicit context dependencies only” [Szyperski 2002]
- Facilitates Reuse of Component Specifications
- Component’s dependence upon entity types is a context dependency
- Approach: Reuse **PCM Resource Interfaces** [Hauck et al. 2009]

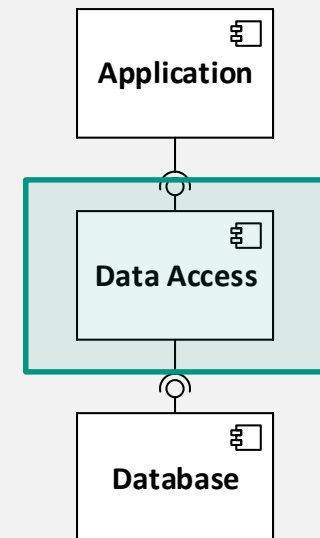
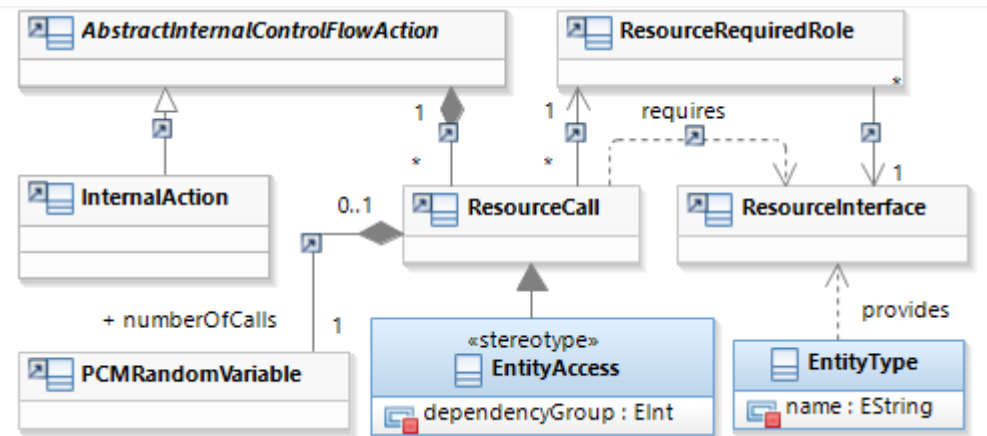


# Schema Definition & Table Deployment



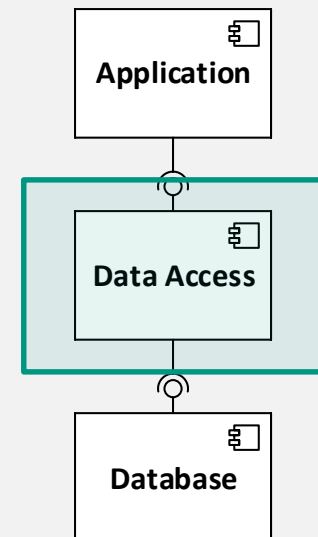
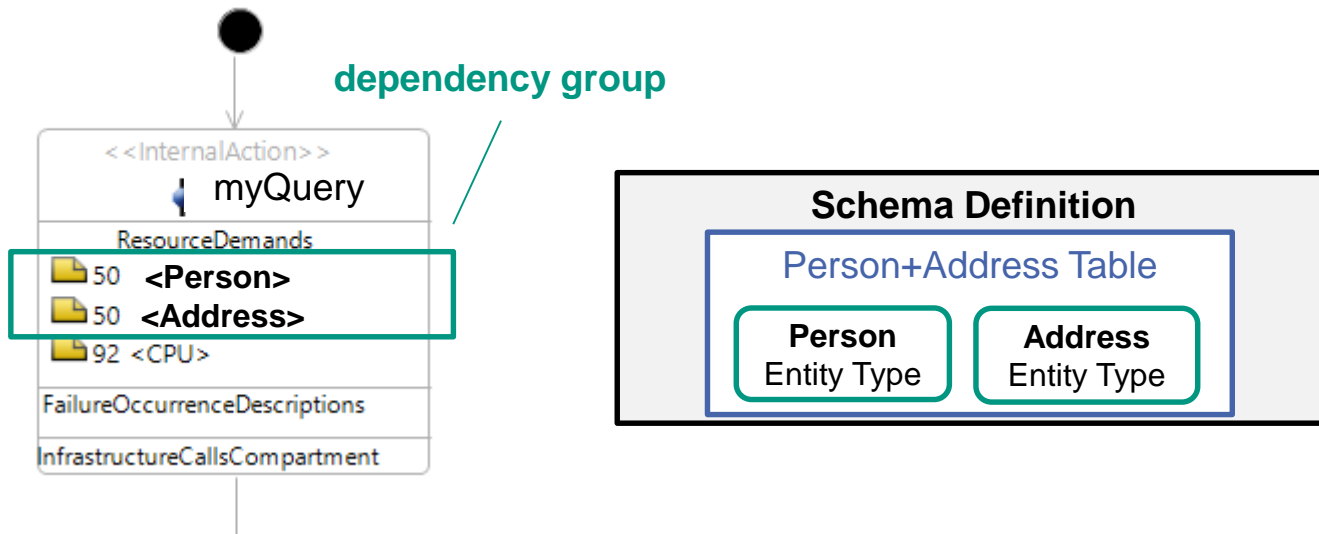
# Queries

- So far: InternalActions require processing resources (CPU, HDD)
- PCM.TX: InternalActions may additionally **require entities**



- *Number of entities* described as PCMRandomVariable, hence a...
  - Constant
  - Probability distribution function
  - Function of (call/component) parameters

# Queries – Dependency Groups



## ■ How many rows from the Person+Address Table are accessed?

- Case 1: 50 rows, if each address **belongs to** one of the 50 persons

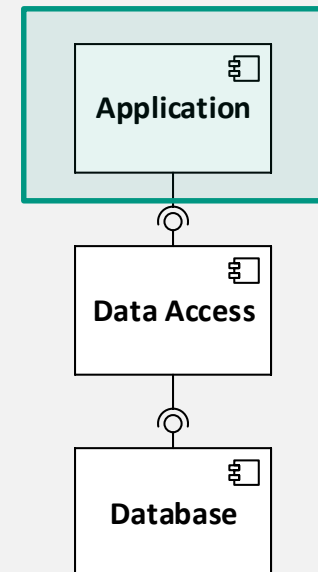
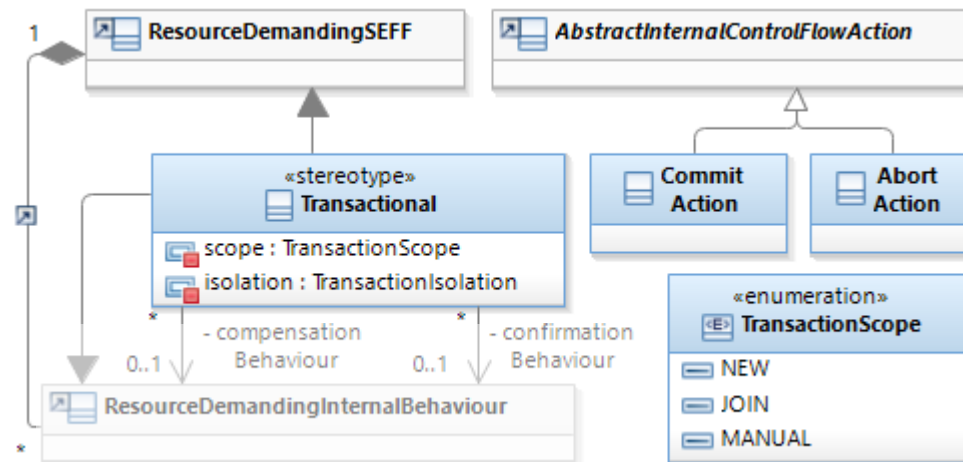
→ assign entity accesses to same dependency group

- Case 2: 100 rows, if there is **no correspondence** between the addresses and persons

→ no dependency (default case)

# Transaction Boundaries

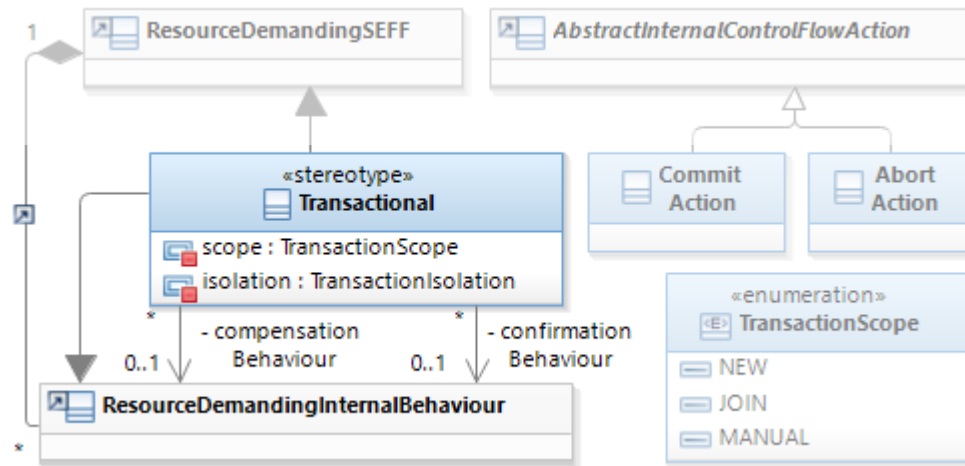
- Annotate transactional SEFFs („declarative“ approach)
- Scope determines what happens when SEFF S1 calls S2
  - Case 1: Two separate transactions, one for S1, one for S2 (NEW)
  - Case 2: A joint transaction (JOIN)



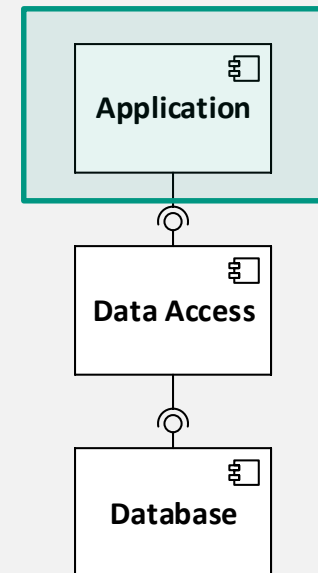
- Scope MANUAL for fine-grained transaction demarcation, together with Commit/Abort

# Compensation, Try-Cancel-Confirm

- Compensation behaviour executes when its corresponding transaction **aborts**



- Confirmation behaviour executes when its corresponding transaction **succeeds**

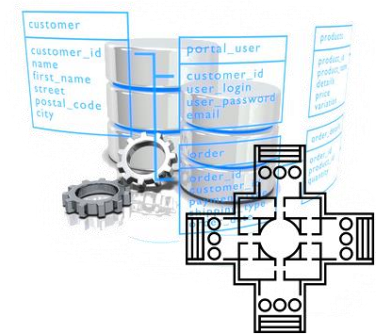


# Conclusion and Future Work

- PCM.TX adds transaction and database modeling to the PCM
- Allows to evaluate design decisions specific to transactional information systems

→ **evaluate impact of migration towards service-orientation**

- Future Work
  - Simulation of PCM.TX models in EventSim
  - Non-uniformly distributed data accesses
  - Sharding (horizontal partitioning of tables)



# References

- **[Szyperski 2002]** C. Szyperski, D. Gruntz, and S. Murer. *Component Software: Beyond Object-Oriented Programming*. ACM Press and Addison-Wesley, New York, 2 edition, 2002.
- **[Hauck et al. 2009]** M. Hauck, M. Kuperberg, K. Krogmann, and R. Reussner. Modelling Layered Component Execution Environments for Performance Prediction. In Proceedings of the 12th International Symposium on Component Based Software Engineering (CBSE 2009), number 5582 in LNCS, pages 191{208. Springer, 2009.



## Backup Slides

# PCM Resource Interfaces

